

Vapnik-Chervonenkis Density in Model Theory

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Let (X, \mathcal{S}) be a set system, i.e., X is a set (the **base set**), and \mathcal{S} is a collection of subsets of X . (We sometimes also speak of a **set system \mathcal{S} on X** .)

Given $A \subseteq X$, we let

$$\mathcal{S} \cap A := \{S \cap A : S \in \mathcal{S}\}$$

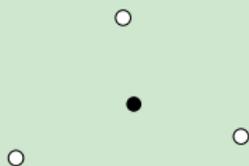
and call $(A, \mathcal{S} \cap A)$ **the set system on A induced by \mathcal{S}** .

We say A is **shattered by \mathcal{S}** if $\mathcal{S} \cap A = 2^A$.

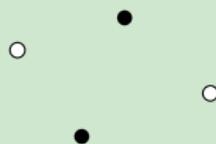
If $\mathcal{S} \neq \emptyset$, then we define the **VC dimension of \mathcal{S}** , denoted by $\text{VC}(\mathcal{S})$, as the supremum (in $\mathbb{N} \cup \{\infty\}$) of the sizes of all finite subsets of X shattered by \mathcal{S} . We also decree $\text{VC}(\emptyset) := -\infty$.

Examples

- 1 $X = \mathbb{R}$, $\mathcal{S} =$ all unbounded intervals. Then $VC(\mathcal{S}) = 2$.
- 2 $X = \mathbb{R}^2$, $\mathcal{S} =$ all halfspaces. Then $VC(\mathcal{S}) = 3$.



One point in the convex hull
of the others

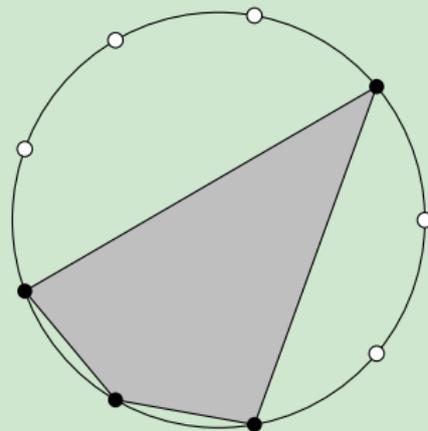
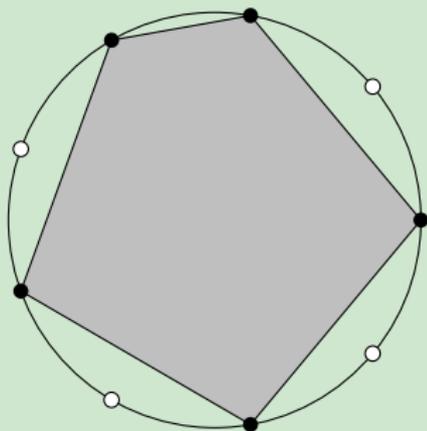


No point in the convex hull
of the others

- 3 Let $\mathcal{S} =$ half spaces in \mathbb{R}^d . Then $VC(\mathcal{S}) = d + 1$.
(The inequality \leq follows from *Radon's Lemma*.)

Examples (continued)

- ④ $X = \mathbb{R}^2$, $\mathcal{S} = \text{all convex polygons}$. Then $\text{VC}(\mathcal{S}) = \infty$.



(But $\text{VC}(\{\text{convex } n\text{-gons in } \mathbb{R}^2\}) = 2n + 1$.)

The function

$$n \mapsto \pi_{\mathcal{S}}(n) := \max \left\{ |\mathcal{S} \cap A| : A \in \binom{X}{n} \right\} : \mathbb{N} \rightarrow \mathbb{N}$$

is called the **shatter function of \mathcal{S}** . Then

$$\text{VC}(\mathcal{S}) = \sup \{n : \pi_{\mathcal{S}}(n) = 2^n\}.$$

One says that \mathcal{S} is a **VC class** if $\text{VC}(\mathcal{S}) < \infty$.

The notion of VC dimension was introduced by Vladimir Vapnik and Alexey Chervonenkis in the early 1970s, in the context of computational learning theory.



A surprising dichotomy holds for $\pi_{\mathcal{S}}$:

The Sauer-Shelah dichotomy

Either

- $\pi_{\mathcal{S}}(n) = 2^n$ for every n (if \mathcal{S} is not a VC class),

or

- $\pi_{\mathcal{S}}(n) \leq \binom{n}{\leq d} := \binom{n}{0} + \dots + \binom{n}{d}$ where $d = \text{VC}(\mathcal{S}) < \infty$.

One may now define the **VC density** of \mathcal{S} as

$$\text{vc}(\mathcal{S}) = \begin{cases} \inf\{r \in \mathbb{R}^{>0} : \pi_{\mathcal{S}}(n) = O(n^r)\} & \text{if } \text{VC}(\mathcal{S}) < \infty \\ \infty & \text{otherwise.} \end{cases}$$

We also define $\text{vc}(\emptyset) := -\infty$.

Examples

- 1 $\mathcal{S} = \binom{X}{\leq d}$. Then $VC(\mathcal{S}) = vc(\mathcal{S}) = d$; in fact $\pi_{\mathcal{S}}(n) = \binom{n}{\leq d}$.
- 2 $\mathcal{S} =$ half spaces in \mathbb{R}^d . Then $VC(\mathcal{S}) = d + 1$, $vc(\mathcal{S}) = d$.

VC density is often the right measure for the combinatorial complexity of a set system.

Some basic properties:

- $vc(\mathcal{S}) \leq VC(\mathcal{S})$, and if one is finite then so is the other;
- $VC(\mathcal{S}) = 0 \iff |\mathcal{S}| = 1$;
- \mathcal{S} is finite $\iff vc(\mathcal{S}) = 0 \iff vc(\mathcal{S}) < 1$;
- $\mathcal{S} = \mathcal{S}_1 \cup \mathcal{S}_2 \Rightarrow vc(\mathcal{S}) = \max\{vc(\mathcal{S}_1), vc(\mathcal{S}_2)\}$.

Let X be a set (possibly finite). Given $A_1, \dots, A_n \subseteq X$, denote by $S(A_1, \dots, A_n)$ the set of atoms of the Boolean subalgebra of 2^X generated by A_1, \dots, A_n : those subsets of X of the form

$$\bigcap_{i \in I} A_i \cap \bigcap_{i \notin I} X \setminus A_i \quad \text{where } I \subseteq \{1, \dots, n\}$$

which are *non-empty* (= “the non-empty sets in the Venn diagram of A_1, \dots, A_n ”).

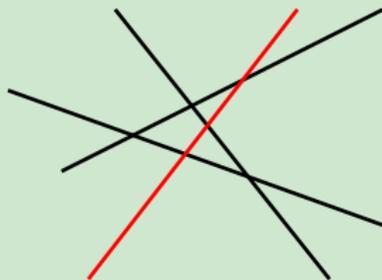
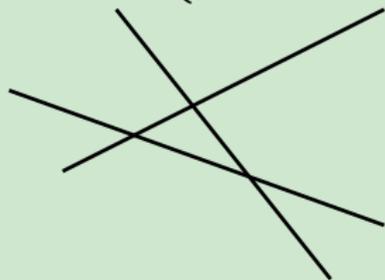
Suppose now that \mathcal{S} is a set system on X . We define

$$n \mapsto \pi_{\mathcal{S}}^*(n) := \max \{ |S(A_1, \dots, A_n)| : A_1, \dots, A_n \in \mathcal{S} \} : \mathbb{N} \rightarrow \mathbb{N}.$$

We say that \mathcal{S} is **independent** (in X) if $\pi_{\mathcal{S}}^*(n) = 2^n$ for every n , and **dependent** (in X) otherwise.

Example ($X = \mathbb{R}^2$, $\mathcal{S} =$ half planes in \mathbb{R}^2)

$\pi_{\mathcal{S}}^*(n) = \left\{ \begin{array}{l} \text{maximum number of regions into which } n \text{ half} \\ \text{planes partition the plane.} \end{array} \right.$



Adding one half plane to $n - 1$ given half planes divides at most n of the existing regions into 2 pieces. So $\pi_{\mathcal{S}}^*(n) = O(n^2)$.

The function $\pi_{\mathcal{S}}^*$ is called the **dual shatter function of \mathcal{S}** .

Let X, Y be infinite sets, $\Phi \subseteq X \times Y$ a binary relation. Put

$$\mathcal{S}_\Phi := \{\Phi_y : y \in Y\} \subseteq 2^X \quad \text{where } \Phi_y := \{x \in X : (x, y) \in \Phi\},$$

and

$$\begin{aligned} \pi_\Phi &:= \pi_{\mathcal{S}_\Phi}, & \pi_\Phi^* &:= \pi_{\mathcal{S}_\Phi^*}, \\ \text{VC}(\Phi) &:= \text{VC}(\mathcal{S}_\Phi), & \text{vc}(\Phi) &:= \text{vc}(\mathcal{S}_\Phi). \end{aligned}$$

We also write

$$\Phi^* \subseteq Y \times X := \{(y, x) \in Y \times X : (x, y) \in \Phi\}.$$

In this way we obtain two set systems: (X, \mathcal{S}_Φ) and $(Y, \mathcal{S}_{\Phi^*})$

Given a finite set $A \subseteq X$ we have a bijection

$$A' \mapsto \bigcap_{x \in A'} \Phi_x^* \cap \bigcap_{x \in A \setminus A'} Y \setminus \Phi_x^* : \mathcal{S}_\Phi \cap A \rightarrow \mathcal{S}(\Phi_x^* : x \in A).$$

Hence $\pi_{\Phi} = \pi_{\Phi^*}^*$ and $\pi_{\Phi^*} = \pi_{\Phi}^*$, and thus

$$\begin{aligned} \mathcal{S}_{\Phi} \text{ is a VC class} &\iff \mathcal{S}_{\Phi^*} \text{ is dependent,} \\ \mathcal{S}_{\Phi^*} \text{ is a VC class} &\iff \mathcal{S}_{\Phi} \text{ is dependent.} \end{aligned}$$

Moreover (first noticed by Assouad):

$$\mathcal{S}_{\Phi} \text{ is a VC class} \iff \mathcal{S}_{\Phi^*} \text{ is a VC class.}$$

We fix:

\mathcal{L} : a first-order language,

$x = (x_1, \dots, x_m)$: object variables,

$y = (y_1, \dots, y_n)$: parameter variables,

$\varphi(x; y)$: a partitioned \mathcal{L} -formula,

M : an infinite \mathcal{L} -structure, and

T : a complete \mathcal{L} -theory without finite models.

The set system (on M^m) associated with φ in M :

$$\mathcal{S}_\varphi^M := \{\varphi^M(M^m; b) : b \in M^n\}$$

If $M \equiv N$, then $\pi_{\mathcal{S}_\varphi^M} = \pi_{\mathcal{S}_\varphi^N}$. So, picking $M \models T$ arbitrary, set

$$\pi_\varphi := \pi_{\mathcal{S}_\varphi^M}, \quad \text{VC}(\varphi) := \text{VC}(\mathcal{S}_\varphi^M), \quad \text{vc}(\varphi) := \text{vc}(\mathcal{S}_\varphi^M).$$

The *dual* of $\varphi(x; y)$ is $\varphi^*(y; x) := \varphi(x; y)$. Put

$$\text{VC}^*(\varphi) := \text{VC}(\varphi^*), \quad \text{vc}^*(\varphi) := \text{vc}(\varphi^*).$$

We have $\pi_\varphi^* = \pi_{\varphi^*}$, hence $\text{VC}^*(\varphi)$ and $\text{vc}^*(\varphi)$ can be computed using the dual shatter function of φ .

If $\text{VC}(\varphi) < \infty$ then we say that φ is **dependent** in T . The theory T does **not have the independence property** (is **NIP**) if every partitioned \mathcal{L} -formula is dependent in T .

An important theorem of Shelah (given other proofs by Laskowski and others) says that for T to be NIP it is enough for every \mathcal{L} -formula $\varphi(x; y)$ with $|x| = 1$ to be dependent.

Many (but not all) well-behaved theories arising naturally in model theory are NIP.

Some questions about vc in model theory

- ① Possible values of $\text{vc}(\varphi)$. There exists a formula $\varphi(x; y)$ in $\mathcal{L}_{\text{rings}}$ with $|y| = 4$ such that

$$\text{vc}^{\text{ACF}_0}(\varphi) = \frac{4}{3}; \quad \text{vc}^{\text{ACF}_p}(\varphi) = \frac{3}{2} \text{ for } p > 0.$$

We do not know an example of a formula φ in a NIP theory with $\text{vc}(\varphi) \notin \mathbb{Q}$.

- ② Growth of π_φ . There is an example of an ω -stable T and an \mathcal{L} -formula $\varphi(x; y)$ with $|y| = 2$ and

$$\pi_\varphi(n) = \frac{1}{2}n \log n (1 + o(1)).$$

- ③ Uniform bounds on $\text{vc}(\varphi)$.

Some reasons why it should be interesting to obtain bounds on $vc(\varphi)$ in terms of $|y| =$ number of free parameters:

- 1 uniform bounds on VC density often “explain” why certain bounds on the complexity of geometric arrangements, used in computational geometry, are polynomial in the number of objects involved;
- 2 connections to strengthenings of the NIP concept: if $vc(\varphi) < 2$ for each $\varphi(x; y)$ with $|y| = 1$ then T is *dp-minimal*.

Theorem

Suppose T expands the theory of linearly ordered sets, and assume that T is **weakly o-minimal**, i.e., in every $M \models T$, every definable subset of M is a finite union of convex subsets of M . Then for each $\varphi(x; y)$ we have $\pi_\varphi(t) = O(t^{|y|})$, hence $\text{vc}(\varphi) \leq |y|$.

- Generalizes results due to Karpinski-Macintyre and Wilkie;
- Idea of the proof:
 - generalize definition of π_φ^* to finite sets Δ of formulas instead of a single φ ;
 - the number of parameters needed in a *uniform definition of Δ -types* over finite parameter sets yields a bound on $\pi_\Delta^*(t)$;
 - reduce to the case where $|x| = 1$ and each instance of $\varphi \in \Delta$ defines an initial segment, in which case each finite Δ -type can be defined by a single parameter.

Uniform bounds on VC density

Interesting classes of NIP theories are provided by certain valued fields. By a non-trivial elaboration of our methods:

Theorem

Suppose $M = \mathbb{Q}_p$ is the field of p -adic numbers, construed as a structure in the language of rings. Then $\text{vc}(\varphi) \leq 2|y| - 1$.

We also have uniform bounds on VC density (obtained by other techniques) for certain stable structures; e.g., we characterize those abelian groups for which we have such uniform bounds.

There are many open questions in this subject.

Open question

If there is some d_1 such that $\text{vc}(\varphi) \leq d_1$ for each $\varphi(x; y)$ with $|y| = 1$, is there is some d_m such that $\text{vc}(\varphi) \leq d_m$ for each $\varphi(x; y)$ with $|y| = m$?

Question

Let $f: A \rightarrow \mathbb{R}^n$, $A \subseteq \mathbb{R}^m$, be L -Lipschitz (where $L \in \mathbb{R}^{\geq 0}$), i.e.,

$$\|f(x) - f(y)\| \leq L \cdot \|x - y\| \quad \text{for all } x, y \in A.$$

Can one extend f to an L -Lipschitz map $\mathbb{R}^m \rightarrow \mathbb{R}^n$?

Kirszbraun (1934): yes for all n

There always exists an L -Lipschitz extension $\mathbb{R}^m \rightarrow \mathbb{R}^n$ of f .

The usual proofs of this theorem all use some sort of transfinite induction.

Question (Chris Miller)

Let $f: A \rightarrow \mathbb{R}^n$, $A \subseteq \mathbb{R}^m$, be L -Lipschitz and *semialgebraic*. Is there a *semialgebraic* L -Lipschitz map $\mathbb{R}^m \rightarrow \mathbb{R}^n$ extending f ?

More generally, one may ask this for Lipschitz maps definable in an o-minimal expansion $\mathbf{R} = (R, 0, 1, +, \times, <, \dots)$ of a real closed field R , instead of $\overline{\mathbb{R}} = (\mathbb{R}, 0, 1, +, \times, <)$.

Why is this extra generality interesting?

- no “higher-order” Tarski principle for transfer from o-minimal expansions of $\overline{\mathbb{R}}$ to \mathbf{R} ;
- brings out the inherent uniformities in the construction.

In fact, o-minimality even turns out to be an unnecessarily strong assumption.

Theorem A (A.-Fischer, Proc. LMS 2011)

Let $\mathcal{R} = (R, 0, 1, +, \times, <, \dots)$ be a **definably complete** expansion of an ordered field: every non-empty definable subset of R which is bounded from above has a supremum. Then every definable L -Lipschitz map $A \rightarrow R^n$ ($A \subseteq R^m$, $L \in R^{\geq 0}$) has a definable L -Lipschitz extension $R^m \rightarrow R^n$.

The proof of this theorem used convex analysis and is based on a relationship between Lipschitz maps and monotone set-valued maps (Minty; more recently, Bauschke & Wang).

Another crucial ingredient (in the case where $R \neq \mathbb{R}$) is a definable version of a classical theorem of Helly:

Theorem B (A.-Fischer, Proc. LMS 2011)

Let R be a definably complete expansion of an ordered field. Let \mathcal{C} be a definable family of closed bounded *convex* subsets of R^n . Suppose \mathcal{C} is $(n + 1)$ -**consistent**:

$$\bigcap \mathcal{C}' \neq \emptyset \quad \text{for all } \mathcal{C}' \subseteq \mathcal{C} \text{ with } |\mathcal{C}'| \leq n + 1.$$

Then $\bigcap \mathcal{C} \neq \emptyset$.

Our proof of this theorem uses an optimization argument.

S. Starchenko pointed out that in the case of an o-minimal R , our theorem follows from an analysis of the model-theoretic notion of *forking* in o-minimal structures due to A. Dolich.

A subset T of X is called a **transversal** of a set system \mathcal{S} on X if every member of \mathcal{S} contains an element of T .

Theorem (Dolich '04, made explicit by Peterzil & Pillay '07)

Let \mathcal{R} be an o-minimal expansion of a real closed field, and let $\mathcal{C} = \{C_a\}_{a \in A}$ be a definable family of closed and bounded subsets of \mathcal{R}^n parameterized by a subset A of \mathcal{R}^m . If \mathcal{C} is $N(m, n)$ -consistent, where

$$N(m, n) = (1 + 2^m) \cdot (1 + 2^{2^m}) \cdots \quad (n \text{ factors}),$$

then \mathcal{C} has a finite transversal.

Question

Can one do better than the bound $N(m, n)$?

Theorem (Matoušek, 2004)

Let (X, \mathcal{S}) be a set system of finite dual VC density $\text{vc}^(\mathcal{S})$. Suppose \mathcal{S} is d -consistent, where $d > \text{vc}^*(\mathcal{S})$. Assume that X comes equipped with a topology making all sets in \mathcal{S} compact. Then \mathcal{S} has a finite transversal.*

Corollary

Let \mathcal{R} be an o-minimal structure on \mathbb{R} , and let $\mathcal{C} = \{C_a\}_{a \in A}$ be a definable family of compact subsets of \mathbb{R}^n . If \mathcal{C} is $(n + 1)$ -consistent, then \mathcal{C} has a finite transversal.

Proof of Theorem B in the o-minimal case (Starchenko)

Suppose R is o-minimal, and write $\mathcal{C} = \{C_a\}_{a \in A}$.

By Helly's Theorem for finite families, the (definable) family whose members are the intersections of $n + 1$ members of \mathcal{C} is finitely consistent.

Apply Dolich's Theorem to this family to obtain a finite set $P \subseteq R^n$ with $P \cap C_{a_1} \cap \cdots \cap C_{a_{n+1}} \neq \emptyset$ for all $a_1, \dots, a_{n+1} \in A$.

Thus

$$\mathcal{P} = \{\text{conv}(C_a \cap P)\}_{a \in A}$$

is a family of convex subsets of R^n with only finitely many distinct members, and \mathcal{P} is $(n + 1)$ -consistent.

Hence $\emptyset \neq \bigcap \mathcal{P} \subseteq \bigcap \mathcal{C}$ by Helly's Theorem for finite families. \square